



SPECIAL RESEARCH REPORT

Reverse engineering of the intrusion behavior
of malware z13.exe

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Introduction

We report here on a reverse engineering effort to better understand how z13.exe infests a system. From our analysis, this malware bears the signature of a banker trojan. Although we have not yet confirmed its insertion in a web browser to steal banking credentials, we can see that it leverages methods and algorithms seen in Zeus and the most recent version of Tinba (Tiny Banker). In addition, it is reported by most antivirus software as miscellaneous banker variations. Since the leak of the Zeus source code in 2011, forks of well-known versions of banker trojans have become commonplace.

We first look at the dropper program, which unwraps the malware behaviour from encrypted modules. Second, we show how it transfers its attack code to a hollowed-out child process. Third, we follow its injection into Windows Explorer, using the so-called Heaven's Gate technique. Fourth, we examine the various tricks the malware does to cloak itself, as well as to persist on the infected system.

Malware Identity

File name / z13.exe

Observed Location /

Architecture / x86 (supports x64 injections)

Arc4dia Indicators

Dropped Files / z13.exe

Autoruns / Key / HKCU\Software\Microsoft\windows\currentVersion\Run

/ Value / C:\%UserDir%\AppData\Roaming\Identities\wnisxpeo.exe lorem ipsum dolor

Architecture / C:\%UserDir%\AppData\Roaming\Identities\wnisxpeo.exe

Table Lorem ipsum

Entry Point

At first glance, the malware doesn't reveal much with static and dynamic analysis. Only with careful tracing does it reveal all its tricks. It starts as a common MFC program. It sets up a dialog named *chat tool* from a loaded resource segment. It's from the initialization of this dialog that the code [Link starts to act like a malware](#).

It loads in memory a copy of the file on disk and then seeks a value into this memory that indicates the start of an encrypted block. The block is decrypted on the heap as follows:

004012DC	B8 OF 82 AE 34	mov eax,34AE820F
004012E1	F7 EE	imul esi
004012E3	C1 FA 06	sar edx,6
004012E6	8B C2	mov eax,edx
004012E8	81 C6 6E 02 00 00	add esi,26E
004012EE	C1 E8 1F	shr eax,1F
004012F1	03 D0	add edx,eax
004012F3	03 D7	add edx,edi
004012F5	8A 85 48 E8 FF FF	mov al,byte ptr ss:[ebp-17B8]
004012F8	8A 94 0A B3 0C 00 00	mov dl,byte ptr ds:[edx+ecx+CB3]
00401302	80 F2 FB	xor dl,FB
00401305	32 D0	xor dl,al
00401307	8B 13	mov byte ptr ds:[ebx],dl
00401309	43	inc ebx
0040130A	81 FE EE 57 1B 00	cmp esi,1B57EE
00401310	0F 8E 25 FF FF FF	jng z13.40123B

Block decryption

Once decrypted, the execution jumps into this new code segment, which executes various recon tasks to determine whether the host is already infected. At the end of this code segment, another memory block is decrypted using the same procedure. This second block contains antiVM behaviour and an interesting forking process detailed in the next section. This type of architecture shows modularity in the conception of the malware. The developers make independent block with different features and those blocks, revealed for execution one at a time, determine the malware behaviour.

02800A45	6A 14	push 14	
02800A47	8D 85 34 FF FF FF	lea eax,dword ptr ss:[ebp-CC]	
02800A4D	57	push edi	
02800A4E	50	push eax	
02800A4F	8D 85 7C C0 FF FF	lea eax,dword ptr ss:[ebp-3F84]	
02800A55	50	push eax	
02800A56	E8 00 F9 FF FF	call 280035B	Decipher function
02800A58	83 C4 10	add esp,10	
02800A5E	6A 40	push 40	
02800A60	68 00 30 00 00	push 3000	
02800A65	57	push edi	
02800A66	53	push ebx	
02800A67	FF 95 48 FF FF FF	call dword ptr ss:[ebp-B8]	[ebp-B8]:VirtualAlloc
02800A6D	8B F0	mov esi,eax	Move newly allocated space ptr (0x3D0000) in ESI
02800A6F	8D 85 7C C0 FF FF	lea eax,dword ptr ss:[ebp-3F84]	
02800A75	57	push edi	
02800A76	50	push eax	
02800A77	56	push esi	
02800A78	E8 B3 F7 FF FF	call 2800230	Move data to allocated space
02800A7D	83 C4 0C	add esp,C	
02800A80	FF D6	call esi	Jump into new code
02800A82	5F	pop edi	
02800A83	5E	pop esi	
02800A84	33 C0	xor eax,eax	
02800A86	5B	pop ebx	
02800A87	C9	leave	
02800A88	C3	ret	
02800A89	CC	int3	

003D0000	00004000		PRV	ERW--	ERW--
00400000	00001000	z13.exe	IMG	-R---	ERWC-
00401000	00002000	".text"	IMG	ER---	ERWC-
00403000	00001000	".rdata"	IMG	-R---	ERWC-
00404000	00001000	".data"	IMG	-RW--	ERWC-
00405000	00001000	".idata"	IMG	-RW--	ERWC-
00406000	0001E000	".rsrc"	IMG	-R---	ERWC-
00424000	00001000	".reloc"	IMG	-R---	ERWC-

Code modularity

The imports for each code segment are found by first pushing the names of all routines it needs on the stack. Then, the string *kernel32.dll* is retrieved from the process environment block (PEB). The module is found among those loaded, so the code fetches *LoadLibrary* and *GetProcAddress* from its the export table. Those two functions are then used to resolve all other function names that were previously pushed on the stack.

02800740	C6 85 7C FF FF FF 2E	mov byte ptr ss:[ebp-84],2E	
02800747	C6 85 7D FF FF FF 65	mov byte ptr ss:[ebp-83],65	
0280074E	C6 85 7E FF FF FF 78	mov byte ptr ss:[ebp-82],78	
02800755	C6 85 7F FF FF FF 65	mov byte ptr ss:[ebp-81],65	
0280075C	C6 45 80 2E	mov byte ptr ss:[ebp-80],2E	
02800760	C6 45 81 5C	mov byte ptr ss:[ebp-7F],5C	
02800764	88 5D 82	mov byte ptr ss:[ebp-7E],b1	
02800767	E8 EC FA FF FF	call 2800258	Load GetProcAddress and LoadLibrary
0280076C	59	pop ecx	
0280076D	8D 85 60 FF FF FF	lea eax,dword ptr ss:[ebp-A0]	
02800773	59	pop ecx	
02800774	50	push eax	
02800775	8D 45 E0	lea eax,dword ptr ss:[ebp-20]	
02800778	50	push eax	
02800779	FF 55 FC	call dword ptr ss:[ebp-4]	[ebp-4]:LoadLibraryA
0280077C	50	push eax	
0280077D	FF 55 F8	call dword ptr ss:[ebp-8]	[ebp-8]:GetProcAddress
02800780	89 85 28 FF FF FF	mov dword ptr ss:[ebp-D8],eax	[ebp-D8]:eFileNameW
02800786	8D 45 B8	lea eax,dword ptr ss:[ebp-48]	
02800789	50	push eax	
0280078A	8D 45 E0	lea eax,dword ptr ss:[ebp-20]	
0280078D	50	push eax	
0280078E	FF 55 FC	call dword ptr ss:[ebp-4]	[ebp-4]:LoadLibraryA
02800791	50	push eax	
02800792	FF 55 F8	call dword ptr ss:[ebp-8]	[ebp-8]:GetProcAddress
02800795	89 45 F0	mov dword ptr ss:[ebp-10],eax	
02800798	8D 45 AC	lea eax,dword ptr ss:[ebp-54]	
0280079B	50	push eax	
0280079C	8D 45 E0	lea eax,dword ptr ss:[ebp-20]	
0280079F	50	push eax	
028007A0	FF 55 FC	call dword ptr ss:[ebp-4]	[ebp-4]:LoadLibraryA
028007A3	50	push eax	
028007A4	FF 55 F8	call dword ptr ss:[ebp-8]	[ebp-8]:GetProcAddress

Import resolving

Malicious behaviour transfer

The malware is using a wellknown but interesting technique called process hollowing, which consists in rewriting the memory of another process without touching the PEB. This technique is usually used to cloak a malicious process into a legitimate process. For example, "C:\malware.exe" could hide into "C:\windows\notepad.exe", but in this case, the hollowed process is the same as the original one. It is therefore not used for cloaking but as a way to slow down the reversing.

push eax	
push ebx	
call dword ptr ss:[ebp-588]	[ebp-588]:CreateProcessW
test eax,eax	
je 39233D	
lea eax,dword ptr ss:[ebp-10C0]	
mov dword ptr ss:[ebp-10C0],10007	
push eax	
push dword ptr ss:[ebp-28]	
call dword ptr ss:[ebp-558]	[ebp-558]:GetThreadContext
push ebx	
lea eax,dword ptr ss:[ebp-1C]	[ebp-1C]:RegQueryValueExA
push 4	
push eax	
push dword ptr ss:[ebp-1010]	
push dword ptr ss:[ebp-2C]	
call dword ptr ss:[ebp-574]	[ebp-574]:ZwReadVirtualMemory
push dword ptr ss:[ebp-1C]	[ebp-1C]:RegQueryValueExA
push dword ptr ss:[ebp-2C]	
call dword ptr ss:[ebp-80]	[ebp-80]:NtUnmapViewOfSection
push dword ptr ss:[ebp-1C]	[ebp-1C]:RegQueryValueExA
push dword ptr ss:[ebp-2C]	
call dword ptr ss:[ebp-80]	[ebp-80]:NtUnmapViewOfSection
push dword ptr ds:[edi+34]	
push dword ptr ss:[ebp-2C]	
call dword ptr ss:[ebp-80]	[ebp-80]:NtUnmapViewOfSection
push dword ptr ds:[edi+34]	
push dword ptr ss:[ebp-2C]	
call dword ptr ss:[ebp-80]	[ebp-80]:NtUnmapViewOfSection
push 40	
push 3000	
push dword ptr ds:[edi+50]	
push dword ptr ds:[edi+34]	
push dword ptr ss:[ebp-2C]	
call dword ptr ss:[ebp-570]	[ebp-570]:VirtualAllocEx
push ebx	

```

push 4
push eax
mov eax,dword ptr ss:[ebp-101C]
add eax,8
push eax
push dword ptr ss:[ebp-2C]
call dword ptr ss:[ebp-4AC]
mov eax,dword ptr ds:[edi+28]
add eax,dword ptr ss:[ebp-48]
mov dword ptr ss:[ebp-1008],eax
mov dword ptr ss:[ebp-1010],eax
lea eax,dword ptr ss:[ebp-10C0]
push eax
push dword ptr ss:[ebp-28]
call dword ptr ss:[ebp-55C]
push dword ptr ss:[ebp-28]
call dword ptr ss:[ebp-560]
cmp dword ptr ss:[ebp+8],7
jnz 392347

```

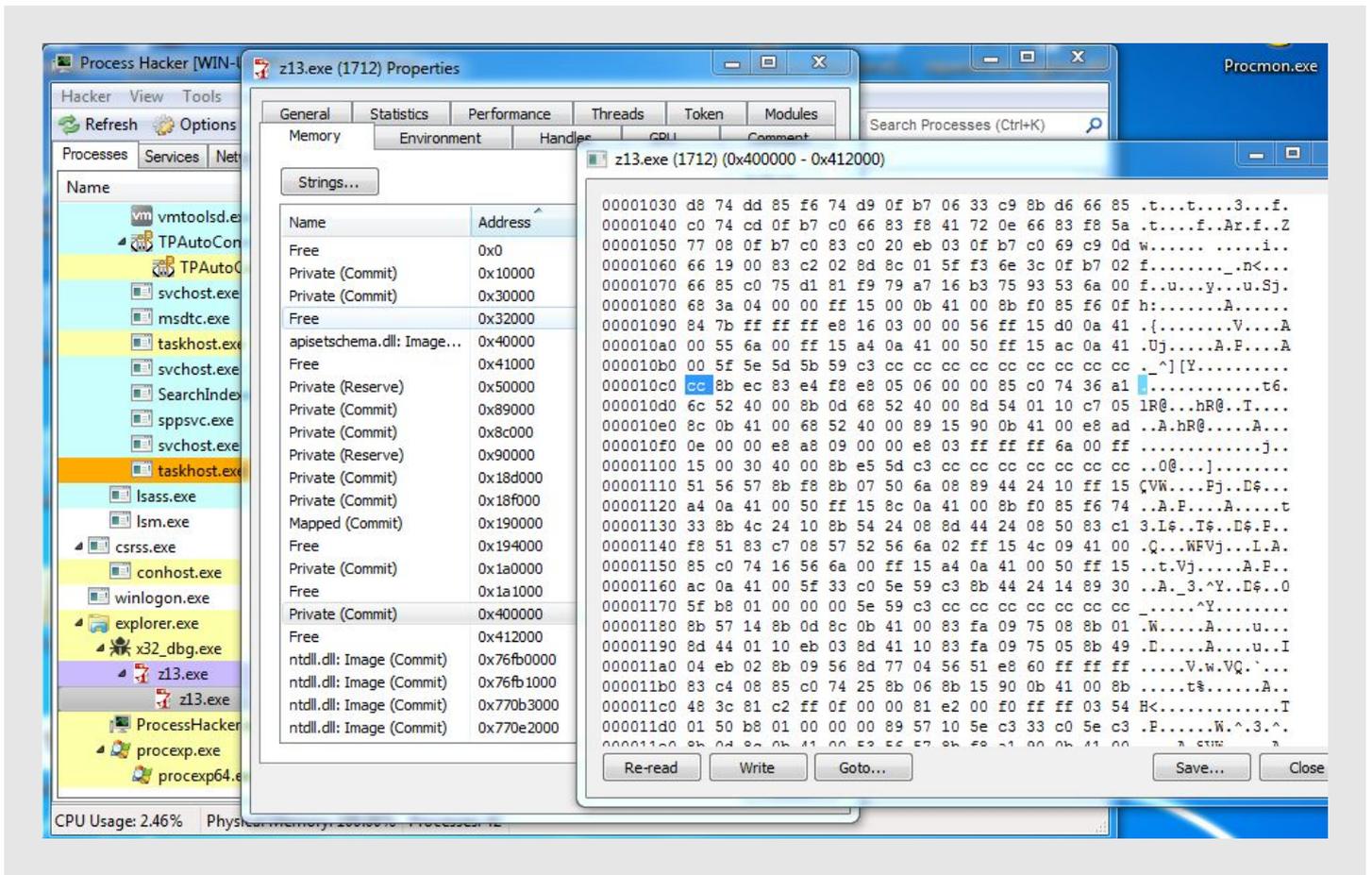
[ebp-4AC]:WriteProcessMemory
[ebp-48]:RegOpenKeyEXA
[ebp-55C]:SetThreadContext
[ebp-560]:ResumeThread

Process hollowing implementation

The implementation of process hollowing starts by creating a suspended target process. The original memory sections of the target process are first unmapped with *NtUnmapViewOfSection* . The new sections are then written into the hollowed out process. The suspended process is now ready to resume except for one detail: the thread EP (entry point) is not valid anymore. Hence the parent process gets the main thread context with *GetThreadContext* , corrects the EP at offset 0xB0 and sets back the thread context with *SetThreadContext* . The hollowed out process is then resumed with new program code.

At this point, the malicious computation has transferred into the child process. In order to continue our trace study, we must attach to the child process. This can be done by manually placing a breakpoint on the new EP with a memory patching tool like ProcessHacker. The EP address is found at offset 0xB0 in the thread context structure just before the *SetThreadContext* function.

Address	Hex	ASCII
00177DDC	07 00 01 00 00 00 00 00 00 00 00 00 00 00 00 00
00177DEC	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00177DFC	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00177E0C	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00177E1C	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00177E2C	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00177E3C	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00177E4C	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00177E5C	00 00 00 00 00 00 00 00 00 00 00 00 00 2B 00 00+...
00177E6C	53 00 00 00 2B 00 00 00 2B 00 00 00 00 00 00 00	S...+...+...
00177E7C	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00áy~.....
00177E8C	C0 10 40 00 00 00 00 00 C0 10 40 00 23 00 00 00	À.@.....À.@.#...
00177E9C	02 02 00 00 F0 FF 18 00 2B 00 00 00 41 50 49 2E	...òÿ...+...API.
00177EAC	64 6C 6C 00 00 00 00 00 00 00 00 00 00 00 00 00	d11.....
00177EBC	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00



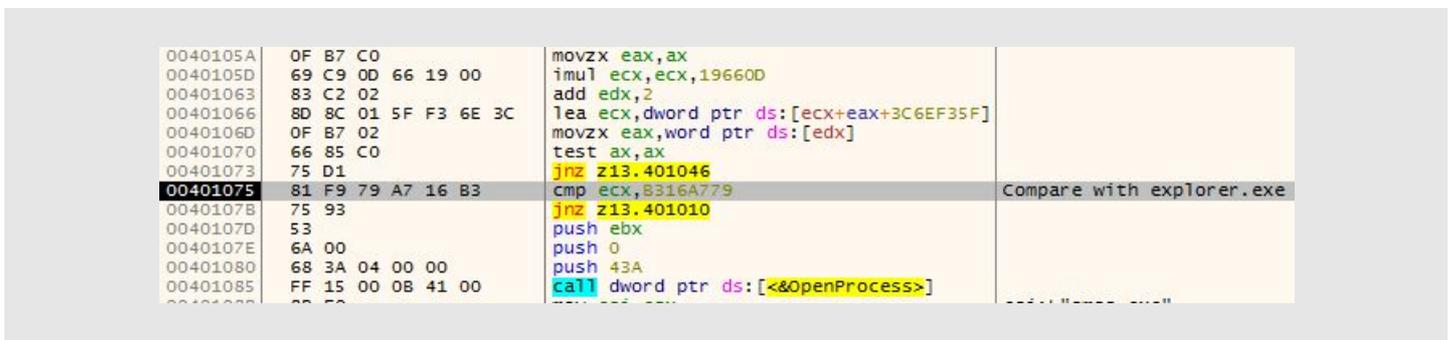
Hollowed process EP

With the breakpoint set in the child process, we set our debugger up as JIT and resume execution of the parent process. The latter will resume the hollowedout process, so the JIT debugger will break on the EP.

The main task of the child process is to inject a remote thread into the Windows Explorer. To prepare for this, it simply asserts debugging privileges, so it may attach as a debugger to all application running as the same user.

Injection in Windows Explorer

The malware finds its injection target by listing the system processes with the function *NtQuerySystemInformation*, called with the enum *SYSTEM_INFORMATION_CLASS* set to *0x5* (*SystemProcessInformation*). It then performs the same hash comparison procedure used to resolve the imports in order to find the *SYSTEM_PROCESS_INFORMATION* structure related to explorer.exe. It gets from that structure the PID of explorer.exe, of which it retrieves the handle using *OpenProcess*.



Explorer.exe hash comparison

When ready to inject, the malware tests the remote process architecture with *isWow64Process*. The injection occurs either ways, but with the difference that it passes through the **Heaven's gate** for a 64bits process.

The Heaven's gate is a simple far jump with the code selector register (CS) set to *0x33*, indicating x64 instructions. Its a simple CPU switch between 32bit and 64bit mode. In *sysWoW64*, this switch is implemented in the function *wow64cpu!X86SwitchTo64BitMode* FS:[0xC0]. In order to keep control over the return address, the injector implemented a widget with a call to offset *0x0* and setting up the return address *0x5* higher on the stack, yielding a return following immediately this fake procedure, into the rest of the injected code..

00402143	6A 33	push 33	CS
00402145	E8 00 00 00 00	call z13.40214A	
0040214A	83 04 24 05	add dword ptr ss:[esp],5	Return address
0040214E	CB	retf	Far return into x64 instructions
0040214F	B8 60 00 00 00	mov eax,60	Start of x64 instructions
00402154	67 65 4C	dec esp	
00402157	8B 10	mov edx,dword ptr ds:[eax]	
00402159	4D	dec ebp	
0040215A	8B 52 18	mov edx,dword ptr ds:[edx+18]	
0040215D	4D	dec ebp	
0040215E	8B 52 10	mov edx,dword ptr ds:[edx+10]	[edx+10]:\t_'B
00402161	49	dec ecx	
00402162	8B 42 30	mov eax,dword ptr ds:[edx+30]	
00402165	48	dec eax	
00402166	85 C0	test eax,eax	
00402168	74 43	je z13.4021AD	
0040216A	4D	dec ebp	

Heaven's gate entrance

Once in 64bits mode, the injection is completed with *CreateRemoteThread*. The inverse mechanism is used to return to x86 mode, setting the CS to 0x23.

004021A8	8B 52 34	mov edx,dword ptr ds:[edx+34]	
004021AB	EB 04	jmp z13.4021B1	
004021AD	33 C0	xor eax,eax	
004021AF	33 D2	xor edx,edx	
004021B1	E8 00 00 00 00	call z13.4021B6	
004021B6	C7 44 24 04 23 00 00 00	mov dword ptr ss:[esp+4],23	CS set back to 0x23
004021BE	83 04 24 0D	add dword ptr ss:[esp],D	
004021C2	CB	retf	Return into x86 mode
004021C3	5D	pop ebp	
004021C4	C3	ret	
004021C5	CC	int3	
004021C6	CC	int3	

Heavens gate exit

The use of this hybrid type of malware is increasingly used by malware authors to adapt to the dual architecture Windows ecosystem. The main thread injected into explorer.exe is behaving as we would expect from any banker trojan. It starts by hooking *ZwResumeThread*, which is called each time a new process is created by explorer.exe.

0000000077001828	0F 05	syscall	
000000007700182A	C3	ret	
000000007700182B	0F 1F 44 00 00	nop dword ptr ds:[rax+rax]	
0000000077001830	E9 23 E9 FE BF	jmp 36FF0158	Hook on ZwResumeThread
0000000077001835	CC	int3	
0000000077001836	CC	int3	
0000000077001837	CC	int3	
0000000077001838	0F 05	syscall	
000000007700183A	C3	ret	
000000007700183B	0F 1F 44 00 00	nop dword ptr ds:[rax+rax]	
0000000077001840	4C 8B D1	mov r10.rcx	

ntdll.ZwResumeThread hooking

The added hook is straightforward. It injects code into each new process if the access rights allow it. Note that the injection still supports both 32bits and 64bits architectures. The main thread injects also into every currently running process.

The injected thread in explorer.exe also tries to beacon its C&C using the new popular thing in the banker world, DGA (domain generator algorithm), with the difference that it generates IP addresses instead of domain names. The malware owners set up a large range of possible IP addresses for their C&C. The DGA queries thousands of IPs until it receives a valid answer from an active C&C. This method of communication is now very popular, as it is very effective against C&C takedowns.

Persistence and cloaking

The malware is using common sandbox evasion behaviours. It searches indicators specific to virtual environments like vmtoolsd.exe (VMWare), VBoxService.exe (VirtualBox) or SbieDll.dll (Sandboxie). It does so by iterating over the module and process list using *CreateToolhelp32Snapshot*. This is easily worked around with memory patching or a VM setup without virtualisation tools.

Before infecting the host, the mutex "qazwsx" is tested to avoid reinfection. It also attempts to open the file "C:\myapp.exe" with *CreateFile*. If the returned handle is valid, the process is terminated without infecting the host. This could be a way for the authors to easily avoid infecting the host in the development process. It tests the string "C:\windows\explorer.exe.\" similarly.

Persistence is achieved by dropping a copy of the malware in the AppData folder. The name of this new file is an obfuscated value that terminates in *wnisxpeo.exe*. Here is the deobfuscation algorithm:

00401C2F	B8 1F 85 EB 51	mov eax,51EB851F
00401C34	F7 E6	mul esi
00401C36	8B CA	mov ecx,edx
00401C38	B8 A3 88 2E BA	mov eax,BA2E88A3
00401C3D	F7 E6	mul esi
00401C3F	8B C2	mov eax,edx
00401C41	C1 E8 05	shr eax,5
00401C44	33 D2	xor edx,edx
00401C46	BF 19 00 00 00	mov edi,19
00401C48	F7 F7	div edi
00401C4D	8B C6	mov eax,esi
00401C4F	C1 E8 05	shr eax,5
00401C52	C1 E9 03	shr ecx,3
00401C55	83 C2 61	add edx,61
00401C58	52	push edx
00401C59	33 D2	xor edx,edx
00401C5B	F7 F7	div edi
00401C5D	B8 25 49 92 24	mov eax,24924925
00401C62	83 C2 61	add edx,61
00401C65	52	push edx
00401C66	F7 E6	mul esi
00401C68	8B C6	mov eax,esi
00401C6A	2B C2	sub eax,edx
00401C6C	D1 E8	shr eax,1
00401C6E	03 C2	add eax,edx
00401C70	C1 E8 04	shr eax,4
00401C73	33 D2	xor edx,edx
00401C75	F7 F7	div edi
00401C77	8B C1	mov eax,ecx
00401C79	6B C9 19	imul ecx,ecx,19
00401C7C	83 C2 61	add edx,61
00401C7F	52	push edx
00401C80	33 D2	xor edx,edx
00401C82	F7 F7	div edi
00401C84	B8 C9 42 16 B2	mov eax,B21642C9
00401C89	83 C2 61	add edx,61
00401C8C	52	push edx
00401C8D	F7 E6	mul esi
00401C8F	8B C2	mov eax,edx

Deobfuscation algorithm for wnixspeo.exe

Address	Hex	ASCII
0018F708	43 00 3A 00 5C 00 55 00 73 00 65 00 72 00 73 00	C.:.\.U.s.e.r.s.
0018F718	5C 00 46 00 6F 00 6F 00 62 00 61 00 72 00 5C 00	\.F.o.o.b.a.r.\.
0018F728	41 00 70 00 70 00 44 00 61 00 74 00 61 00 5C 00	A.p.p.D.a.t.a.\.
0018F738	52 00 6F 00 61 00 6D 00 69 00 6E 00 67 00 5C 00	R.o.o.a.m.i.n.g.\.
0018F748	49 00 64 00 65 00 6E 00 74 00 69 00 74 00 69 00	I.d.e.n.t.i.t.i.
0018F758	65 00 73 00 5C 00 77 00 6E 00 69 00 73 00 78 00	e.s.\.w.n.i.s.x.
0018F768	70 00 65 00 6F 00 2E 00 65 00 78 00 65 00 00 00	p.e.o...e.x.e...
0018F778	D0 34 35 75 44 00 00 00 20 02 00 00 02 00 00 00	D45UD... ..
0018F788	A4 F7 18 00 C9 14 30 75 40 35 35 75 08 00 00 00	#+..É.Ou@55u... ..
0018F798	30 00 00 00 84 F7 18 00 01 00 00 00 AC F8 18 00	Q...+.....~0... ..
0018F7A8	C8 3E 51 00 00 32 35 75 87 0E 30 75 DC F7 18 00	È>Q..25u..OuÛ÷.. ..
0018F7B8	99 49 2F 75 00 00 2E 75 01 00 00 00 00 00 00 00	.I/u...u... ..
0018F7C8	AC F8 18 00 01 00 00 00 F0 F7 18 00 00 00 00 00	~0... .. ð+... ..
0018F7D8	88 F8 00 01 FC F7 18 00 30 99 1C 77 00 00 2E 75	.0..Û÷...0..w...u ..
0018F7E8	01 00 00 00 00 00 00 00 01 00 00 00 AC F8 18 00~0... ..
0018F7F8	C8 3E 51 00 88 F8 18 00 A9 D8 1C 77 E5 49 2F 75	È>Q..0...@0.wâI/u ..

The hiding path

The file time of the dropped file wnixspeo.exe is copied from svchost.exe. This is a common approach to lower the file suspicion.

6A 03	push 3	
6A 00	push 0	
6A 01	push 1	
68 80 00 00 00	push 80	
8D 4C 24 30	lea ecx,dword ptr ss:[esp+30]	
51	push ecx	ecx:L"C:\\windows\\system32\\svchost.exe"
FF 15 BC 0A 41 00	call dword ptr ds:[<&CreateFile>]	
8B F0	mov esi,eax	
83 FE FF	cmp esi,FFFFFFFF	
74 2A	je 213.401D8D	
8D 54 24 10	lea edx,dword ptr ss:[esp+10]	
52	push edx	
6A 00	push 0	
8D 44 24 10	lea eax,dword ptr ss:[esp+10]	
50	push eax	
56	push esi	
FF 15 74 08 41 00	call dword ptr ds:[<&GetFileTime>]	Get svchost.exe file time
85 C0	test eax,eax	
74 13	je 213.401D8D	
8D 4C 24 10	lea ecx,dword ptr ss:[esp+10]	
51	push ecx	ecx:L"C:\\windows\\system32\\svchost.exe"
6A 00	push 0	
8D 54 24 10	lea edx,dword ptr ss:[esp+10]	
52	push edx	
57	push edi	
FF 15 78 08 41 00	call dword ptr ds:[<&SetFileTime>]	Set wnixspeo.exe file time
56	push esi	
FF 15 D0 0A 41 00	call dword ptr ds:[<&CloseHandle>]	
57	push edi	

An autorun registry key is created by the injected explorer.exe process to secure persistence after reboot.

ab (Default)	REG_SZ	(value not set)
ab wnisxpeo.exe	REG_SZ	"C:\Users\Foobar\AppData\Roaming\Identities\wnisxpeo.exe"

Autorun key

Finally, the dropper is deleted from disk once its task is done. It does so by creating a simple batch file starting by "ms" and followed by a random numerical value. The script is placed in the path "C:\%UserDir%\AppData\Roaming" and executed.

```
ms3027322.bat
1  :q
2  if not exist "C:\Users\Foobar\Desktop\z13.exe" goto z
3  del /Q /F "C:\Users\Foobar\Desktop\z13.exe"
4  goto q
5  :z
6  del /Q /F "C:\Users\Foobar\AppData\Roaming\MS3027~1.BAT"
```

Auto delete script